

EDF+TBS*

Aperiodici → BASSA Priorità

Periodici → ALTA Priorità

troviamo le deadline più stringenti possibili per gli aperiodici

complessità algoritmo ↑
Tempo rispetta aperiodici ↓

→ Algoritmo! $d_k^* = f_k$ → tempo di fine stimato
 al worst case di una richiesta schedata con EDF
 un deadline d_k con TBS

iteratione

Stimate → $\tilde{r}_k^s = a_k + c_k + I_p(a_k, d_k^s)$

UPPER BOUND
 $s \times f_k$

interferenza data da i processi periodici.

$$d_0^{(0)} = 0$$

$$d_k^{(0)} = \max\{a_k, d_{k-1}^{(0)}\} + \lceil \frac{c_k}{v_s} \rceil$$

$$d_k^{s+1} = \tilde{r}_k^s \quad \text{se } d_k^{s+1} = d_k^s \rightarrow \text{STOP}$$

$$I_p(a_k, d_k^s) = I_a(a_k, d_k^s) + I_f(a_k, d_k^s)$$

$$I_a(a_k, d_k^s) = \sum_{\substack{\text{Attivo in } a_k \\ d_i^s < d_k^s}} c_i$$

PROCESSI PER. ATTIVI

DOVUTO AI PROCESSI PERIODICI FUTURI

$$J_k(a_k, d_k^s) = \sum_i \text{MAX} \left\{ 0, \left\lfloor \frac{d_k^s - \text{NEXT-}\tau_i(a_k)}{\tau_i} \right\rfloor - 1 \right\} C_i$$

$$\text{NEXT-}\tau_i(a_k) = \left\lfloor \frac{a_{k+1}}{\tau_i} \right\rfloor \tau_i$$

$$\tilde{f}_k^s = d_k^s \Rightarrow \tilde{r}_k^s = d_k^s$$

1. (es. 3.5)

	τ_1	τ_2	τ_3	τ_4
ϕ_i	0	0	0	0
C_i	1	1	2	2
τ_i	4	5	11	16

J_1	
a_i	0
C_i	3

$$U_p = 0,757 < 1 \quad U_s = 1 - U_p = 0,243$$

$$d_0^s = 0 \quad d_1^s = \text{MAX} \left\{ 0, 0 \right\} + \left\lfloor \frac{3}{0,757} \right\rfloor = 3$$

$$\text{NEXT-}\tau_1(0) = \left\lfloor \frac{0+1}{4} \right\rfloor \cdot 4 = 4$$

$$\text{NEXT-}\tau_2(0) = \left\lfloor \frac{0+1}{5} \right\rfloor \cdot 5 = 5$$

$$\text{NEXT-}\tau_3(0) = \left\lfloor \frac{0+1}{11} \right\rfloor \cdot 11 = 11$$

$$\text{NEXT-}\tau_4(0) = \left\lfloor \frac{0+1}{16} \right\rfloor \cdot 16 = 16$$

$$J_k(0, d_1^s) = \text{MAX} \left\{ 0, \left\lfloor \frac{3-4}{4} \right\rfloor - 1 \right\} \cdot 1 + \text{MAX} \left\{ 0, \left\lfloor \frac{3-5}{5} \right\rfloor - 1 \right\} \cdot 1 + \text{MAX} \left\{ 0, \left\lfloor \frac{3-11}{11} \right\rfloor - 1 \right\} \cdot 2 + \text{MAX} \left\{ 0, \left\lfloor \frac{3-16}{16} \right\rfloor - 1 \right\} \cdot 2 = 3$$

Mo	Tu	We	Th	Fr	Sa	Su
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$$I_a(0, d_1^0) = 1+1+2 = 4 \Rightarrow I_p(0, d_1^0) = 7$$

x_1, x_2, x_3
 hanno
 4 caline
 < 13

$$\Rightarrow \tilde{f}_1^0 = 0+3+7 = 10$$

↳ $d_1^1 = \tilde{f}_1^0 = 10 \neq 13$ continuiamo \rightarrow calcoliamo \tilde{f}_1^1

$$I_p(0, d_1^1) = 3 \rightarrow I_p(0, \tilde{d}_1^1) = \max \left\{ 0, \left\lceil \frac{10-4}{4} \right\rceil - 1 \right\} \cdot 1 +$$

$$\max \left\{ 0, \left\lceil \frac{10-5}{5} \right\rceil - 1 \right\} \cdot 1 +$$

$$\max \left\{ 0, \left\lceil \frac{10-11}{11} \right\rceil - 1 \right\} \cdot 2 +$$

$$\max \left\{ 0, \left\lceil \frac{10-16}{16} \right\rceil - 1 \right\} \cdot 2 = 1$$

$$I_a(0, 10) = 1+1 = 2$$

$$\tilde{f}_1^1 = 0+3+3 = 6$$

↳ $d_1^2 = 6 \neq 10$

$$I_p(0, 6) = 2 \rightarrow I_p(0, 6) = 0$$

$$I_a(0, 6) = 2$$

$$\tilde{f}_1^2 = 0+3+2 = 5$$

↳ $d_1^3 = 5 \neq 6$

$$I_p(0, 5) = 1 \rightarrow I_p(0, 5) = 0$$

$$I_a(0, 5) = 1$$

$$\tilde{f}_1^3 = 3+1 = 4$$

$$\rightarrow d_1^4 = 4 \neq 5$$



$$I_p(0,4) = 0 \quad \begin{cases} I_p(0,4) = 0 \\ I_a(0,4) = 0 \end{cases}$$

$$\tilde{r}_1^4 = 3$$

possiamo trovare qui

$$\hookrightarrow d_u^5 = 3 \neq 4$$

$$I_p(0,3) = 0 \quad \begin{cases} I_p(0,3) = 0 \\ I_a(0,3) = 0 \end{cases}$$

minima deadline possibile per $C_n = 3$

$$\tilde{r}_1^5 = 3$$

$$\rightarrow d_1^6 = 3 = d_1^5 \Rightarrow d_1 = 3$$

2. (es. 3.6)

	τ_1	τ_2
ϕ_i	0	0
C_i	1	2
T_i	6	9

	T_1	T_2
a_i	1	8
C_i	2	1

→ ordinati secondo istanti di arrivo T_1, T_2

$$U_p = \frac{1+2}{6+9} = 0,39 \quad U_s = 1 - U_p = 0,61$$

$$d_0^0 = 0$$

$$d_1^0 = \text{MAX} \{1, 0\} + \left\lceil \frac{2}{0,61} \right\rceil = 5$$

$$\text{NEXT } \tau_1(1) = \left\lceil \frac{1+1}{6} \right\rceil \cdot 6 = 6$$

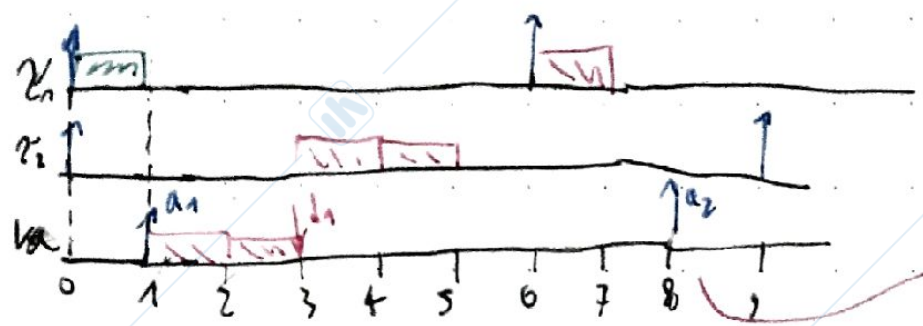
$$\text{NEXT } \tau_2(1) = \left\lceil \frac{1+1}{9} \right\rceil \cdot 9 = 9$$

$$I_p(1,5) = \text{MAX} \{0, \left\lceil \frac{5-6}{6} \right\rceil - 1\} \cdot 1 + \text{MAX} \{0, \left\lceil \frac{5-9}{9} \right\rceil - 1\} \cdot 2 = 0$$

$I_a(1,5)$ → non sappiamo quali sono altri in 1 → schema

Mo	Tu	We	Th	Fr	Sa	Su
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CHECK γ_i $t=1$



nessun processo attivo in $t=3$

γ_2 è attivo, $w_0 > 5$ in $t=1$

$I_e(1, 5) = 0$

$I_p = 0$

$f_1^0 = 1+2 = 3 = a_1 + c_1 \Rightarrow d_1 = 3$

$d_2^0 = \text{MAX} \{ 3, 5 \} + \left\lceil \frac{1}{0.166} \right\rceil = 10$

$\text{NEXT-} \pi_1(2) = \left\lceil \frac{2+1}{6} \right\rceil \cdot 6 = 12$

$I_f(8, 10) = \text{MAX} \{ 0, \left\lceil \frac{10-12}{12} \right\rceil - 1 \} \cdot 1 = 0$

$\text{NEXT-} \pi_2(1) = \left\lceil \frac{1+1}{9} \right\rceil \cdot 9 = 9$

$\text{MAX} \{ 0, \left\lceil \frac{10-9}{9} \right\rceil - 1 \} \cdot 2 = 0$

$\Rightarrow I_p(8, 10) = 0$

$I_a(8, 10) = 0 \rightarrow$ gratuito

$f_2^0 = 8+1 = 9 = a_2 + c_2 \Rightarrow d_2 = 9$